

## Objectives

Reading assignment review  
 Introduction to Algorithms
 

- Our process, through an example

 Representative problems

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## Assignment: Review Paper

Read paper

- 2 hours max
- Don't get bogged down in terminology
- Addendum:
  - Benefits of approach
  - Results of approach

Review paper

- Almost an abstract
- Write on Sakai forum
- Due by 10 a.m. on Friday

Friday: Discuss paper and questions

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## Details

Sakai: Log in

Meeting with Candidates?

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## CS Talks

Attend two CS talks, all in Parmly 405

- Mon, Jan 12, D period
  - Andrea Tartaro: "Authorable Virtual Peers: Using Computer Science to Understand and Support Children with Special Needs"
- Thurs, Jan 15, 3:30 p.m.
  - Mark Liffiton, "Satisfying Constraints, and What To Do When You Can't"
- Fri, Jan 23, 4 p.m.
  - Joshua Stough, "Appearance Models for Medical Image Segmentation"

Post summary on Sakai

- More on that later

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## ALGORITHMS

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## Computational Problem Solving 101

Computational Problem

- A problem that can be solved by logic

To solve the problem:

- Create a **model** of the problem
- Design an **algorithm** for solving the problem using the model
- Write a **program** that *implements* the algorithm

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## Computational Problem Solving 101

Algorithm: a well-defined recipe for solving a problem

- Has a finite number of steps
- Completes in a finite amount of time

Program

- An algorithm written in a **programming language**
- Also called code

Application

- Large programs, solving many problems

## Addendum

Before creating algorithm, define the problem

- Simplify as necessary

## Matching Residents to Hospitals

**Goal:** Given a set of preferences among hospitals and medical school students, design a **self-reinforcing** admissions process.

**Unstable pair:** applicant  $x$  and hospital  $y$  are unstable if:

- $x$  prefers  $y$  to its assigned hospital
- $y$  prefers  $x$  to one of its admitted students

**Stable assignment:** Assignment with no unstable pairs

- Natural and desirable condition
- Individual self-interest will prevent any applicant/hospital deal from being made

What details make this problem tricky?

## Stable Matching Problem

*Simplified version of resident-matching problem*

Goal: Given  $n$  men and  $n$  women, find a "suitable" matching

- Participants rate members of opposite sex
- Each man lists women in order of preference from best to worst
- Each woman lists men in order of preference from best to worst

	favorite ↓	1 <sup>st</sup>	2 <sup>nd</sup>	3 <sup>rd</sup>	least favorite ↓		favorite ↓	1 <sup>st</sup>	2 <sup>nd</sup>	3 <sup>rd</sup>	least favorite ↓
Xavier		Amy	Bertha	Clare		Amy		Yancey	Xavier	Zeus	
Yancey		Bertha	Amy	Clare		Bertha		Xavier	Yancey	Zeus	
Zeus		Amy	Bertha	Clare		Clare		Xavier	Yancey	Zeus	

Men's Preference Profile      Women's Preference Profile

## Stable Matching Problem

**Perfect matching:** everyone is matched monogamously

- Each man is paired with exactly one woman
- Each woman is paired with exactly one man

**Stability:** no incentive for some pair of participants to undermine assignment by joint action

- In matching  $M$ , an *unmatched* pair  $m-w$  is unstable if man  $m$  and woman  $w$  prefer each other to current partners
- Unstable pair  $m-w$  could each improve by eloping

**Stable matching:** perfect matching with no unstable pairs

**Stable matching problem.** Given the preference lists of  $n$  men and  $n$  women, find a stable matching if one exists.

## Analyzing Stability

Q. Is assignment X-C, Y-B, Z-A stable?

- **Unstable:**  $m$  prefers  $w$  to his woman;  $w$  prefers  $m$  to her man

	favorite ↓	1 <sup>st</sup>	2 <sup>nd</sup>	3 <sup>rd</sup>	least favorite ↓		favorite ↓	1 <sup>st</sup>	2 <sup>nd</sup>	3 <sup>rd</sup>	least favorite ↓
Xavier		Amy	Bertha	Clare		Amy		Yancey	Xavier	Zeus	
Yancey		Bertha	Amy	Clare		Bertha		Xavier	Yancey	Zeus	
Zeus		Amy	Bertha	Clare		Clare		Xavier	Yancey	Zeus	

Men's Preference Profile      Women's Preference Profile

### Stable Matching Problem

Q. Is assignment X-C, Y-B, Z-A stable?  
 A. No. Bertha and Xavier prefer each other

	favorite ↓ 1 <sup>st</sup>	2 <sup>nd</sup>	3 <sup>rd</sup>	least favorite ↓		favorite ↓ 1 <sup>st</sup>	2 <sup>nd</sup>	3 <sup>rd</sup>	least favorite ↓
Xavier	Amy	Bertha	Clare		Amy	Yancey	Xavier	Zeus	
Yancey	Bertha	Amy	Clare		Bertha	Xavier	Yancey	Zeus	
Zeus	Amy	Bertha	Clare		Clare	Xavier	Yancey	Zeus	

Men's Preference Profile                      Women's Preference Profile

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### Stable Matching Problem

Q. Is assignment X-A, Y-B, Z-C stable?  
 A. Yes.

	favorite ↓ 1 <sup>st</sup>	2 <sup>nd</sup>	3 <sup>rd</sup>	least favorite ↓		favorite ↓ 1 <sup>st</sup>	2 <sup>nd</sup>	3 <sup>rd</sup>	least favorite ↓
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Yancey	Bertha	Amy	Clare		Bertha	Xavier	Yancey	Zeus	
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### Any Questions?

What are you wondering about this problem at this point?

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### Propose-And-Reject Algorithm

[Gale-Shapley 1962]

Intuitive method that guarantees to find a stable matching

```

Initialize each person to be free.
while (some man is free and hasn't proposed to every woman) {
  Choose such a man m
  w = 1st woman on m's list to whom m has not yet proposed
  if (w is free)
    assign m and w to be engaged
  else if (w prefers m to her fiancé m')
    assign m and w to be engaged, and m' to be free
  else
    w rejects m
}
    
```

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### Applying the Algorithm

	1 <sup>st</sup>	2 <sup>nd</sup>	3 <sup>rd</sup>		1 <sup>st</sup>	2 <sup>nd</sup>	3 <sup>rd</sup>	
Xavier	Amy	Bertha	Clare		Amy	Yancey	Xavier	Zeus
Yancey	Bertha	Amy	Clare		Bertha	Xavier	Yancey	Zeus
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}
    
```

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### Observations about the Algorithm

What can we say about any woman's partner over the execution of the algorithm?

How does a woman's state change over the execution of the algorithm?

What can we say about a man's partner?

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### Algorithm Analysis

What is the running time of this algorithm?  
 What is the state complexity of this algorithm?

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### Proof of Correctness: Termination

**Observation 1.** Men propose to women in decreasing order of preference  
**Observation 2.** Once a woman is matched, she never becomes unmatched; she only "trades up"  
**Claim.** Algorithm terminates after at most  $n^2$  iterations of while loop.

	1 <sup>st</sup>	2 <sup>nd</sup>	3 <sup>rd</sup>	4 <sup>th</sup>	5 <sup>th</sup>		1 <sup>st</sup>	2 <sup>nd</sup>	3 <sup>rd</sup>	4 <sup>th</sup>	5 <sup>th</sup>
Victor	A	B	C	D	E	Amy	W	X	Y	Z	V
Wyatt	B	C	D	A	E	Bertha	X	Y	Z	V	W
Xavier	C	D	A	B	E	Clare	Y	Z	V	W	X
Yancy	D	A	B	C	E	Diane	Z	V	W	X	Y
Zeus	A	B	C	D	E	Enka	V	W	X	Y	Z

$n(n-1) + 1$  proposals required

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### Proof of Correctness: Termination

**Observation 1.** Men propose to women in decreasing order of preference  
**Observation 2.** Once a woman is matched, she never becomes unmatched; she only "trades up"  
**Claim.** Algorithm terminates after at most  $n^2$  iterations of while loop.  
**Pf.** Each time through the while loop a man proposes to a new woman. There are only  $n^2$  possible proposals.

$n(n-1) + 1$  proposals required

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### Assignments

- Read Tartaro paper
  - Discuss on Friday
- Read Chapter 1 of Algorithms book
  - Return to next Friday

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