

Objectives

- Introduction to Algorithms, Analysis
- Course summary
- Reviewing proof techniques

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What This Course Is About



From
30 Rock

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Now, everything comes down to expert knowledge of **algorithms** and **data structures**. If you don't speak fluent **O-notation**, you may have trouble getting your next job at the technology companies in the forefront.

-- Larry Freeman

For me, great algorithms are the poetry of computation. Just like verse, they can be terse, allusive, dense, and even mysterious. But once unlocked, they cast a **brilliant new light** on some aspect of computing.

-- Francis Sullivan

What is an Algorithm?

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Questions to Consider

- What are our goals when designing algorithms?
- How do we know when we've met our goals?

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Course Goals

- Learn how to formulate precise problem descriptions
- Learn specific algorithm design techniques and how to apply them
- Learn how to analyze algorithms for efficiency and for correctness
- Learn when no exact, efficient solution is possible

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Course Content

- Algorithm analysis
 - Formal – proofs; Asymptotic bounds
- Advanced data structures, e.g., heaps, graphs
- Greedy Algorithms
- Dynamic Programming
- Divide and Conquer
- Network Flow
- Computational Intractability

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Course Notes


- Textbook: *Algorithm Design*
 - Optional: CLRS
- Participation is encouraged
 - Individual, group, class
- Assignments:
 - Solutions to problems
 - Analysis of solutions
 - Programming

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Course Grading

- 40% Individual written and programming homework assignments
- 10% Quizzes 
- 25% Midterms
- 20% Final
- 5% Participation and attendance

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ALGORITHMS

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Computational Problem Solving 101

- Computational Problem
 - A problem that can be solved by logic
- To solve the problem:
 1. Create a model of the problem
 2. Design an algorithm for solving the problem using the model
 3. Write a program that implements the algorithm

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Computational Problem Solving 101

- Algorithm: a well-defined recipe for solving a problem
 - Has a finite number of steps
 - Completes in a finite amount of time
- Program
 - An algorithm written in a programming language

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PROOFS

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Why Proofs?

- What are insufficient alternatives?
- How can we prove something isn't true?

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Why Proofs?

- What are insufficient alternatives?
 - Examples
 - Considered all possible?
 - Empirical/statistical evidence
 - Ex: "Lying" with statistics
- How can we prove something isn't true?
 - One counterexample

Need irrefutable proof that something is true—for **all** possibilities

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Common Types of Proofs

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Common Types of Proofs

- Proof by contradiction
- Proof by induction

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Proof By Contradiction

What are the steps to a proof by contradiction?

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Proof By Contradiction

1. Assume the thing we want to prove is false
2. Reason to a contradiction
3. Conclude that it must therefore be true

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Example: There are Infinitely Many Primes

- What is a prime number?
- What is not-a-prime number?

- Assume there are only finitely many prime numbers
 - List them: p_1, p_2, \dots, p_n

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Example: There are Infinitely Many Primes

- Assume there are only finitely many prime numbers
 - List them: p_1, p_2, \dots, p_n
- Consider the number $q = p_1 p_2 \dots p_n + 1$

What are the possibilities for q ?

q is either composite or prime

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Example: There are Infinitely Many Primes

- Assume there are only finitely many prime numbers
 - List them: p_1, p_2, \dots, p_n
- Consider the number $q = p_1 p_2 \dots p_n + 1$
- Case: q is composite
 - If we divide q by any of the primes, we get a remainder of 1 $\rightarrow q$ is not composite

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Example: There are Infinitely Many Primes

- Assume there are only finitely many prime numbers
 - List them: p_1, p_2, \dots, p_n
- Consider the number $q = p_1 p_2 \dots p_n + 1$
- Case: q is composite
 - If we divide q by any of the primes, we get a remainder of 1 $\rightarrow q$ is not composite
- Therefore, q is prime, but q is larger than any of the finitely enumerated prime numbers listed \rightarrow **Contradiction**

Proof thanks
to Euclid

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Proof By Induction

What are the steps to a proof by induction?

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Proof By Induction

1. What you want to prove
2. Base case
 - Typical: Show statement holds for $n = 0$ or $n = 1$
3. Assumption for n (**induction hypothesis**)
4. Induction step: show that adding one to n also holds true
 - Often relies on earlier assumptions

When/why is induction useful?

Show true for all (infinite) possibilities
Show works for "one more"

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Warm Up

Prove:

$$2+4+6+8+\dots + 2n = n*(n+1)$$

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Proof

Prove: $2+4+6+8+\dots + 2n = n*(n+1)$

- **Base case:** $n = 1 \rightarrow 2*1 = 1*(1+1)$ ✓
- Assume true for n
- Prove for $n+1$
 - $2+4+6+8+\dots + 2n + 2(n+1)$
 - $= n*(n+1) + 2(n+1)$
 - $= n^2 + n + 2n + 1 = n^2 + 3n + 1$
 - $= (n+1)*(n+2)$
 - $= (n+1)*((n+1)+1)$ ✓

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Proof: All Horses Are The Same Color

- **Base case:** If there is only *one* horse, there is only one color.
- **Induction step:** Assume as induction hypothesis that within any set of n horses, there is only one color.
 - Look at any set of $n + 1$ horses
 - Label the horses: $1, 2, 3, \dots, n, n + 1$
 - Consider the sets $\{1, 2, 3, \dots, n\}$ and $\{2, 3, 4, \dots, n + 1\}$
 - Each is a set of only n horses, therefore within each there is only one color
 - Since the two sets overlap, there must be only one color among all $n + 1$ horses

Where is the error in the proof?

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Error in Proof

- **Base case:** If there is only *one* horse, there is only one color.
- **Induction step:** Assume as induction hypothesis that within any set of n horses, there is only one color.
 - Look at any set of $n + 1$ horses
 - Number them: $1, 2, 3, \dots, n, n + 1$
 - Consider the sets $\{1, 2, 3, \dots, n\}$ and $\{2, 3, 4, \dots, n + 1\}$
 - Each is a set of only n horses, therefore within each there is only one color
 - Since the two sets overlap, there must be only one color among all $n + 1$ horses

Does not hold true when $n+1=2$

Lesson: check assumptions within proof

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Proof Summary

- Need to prove conjectures
- Common types of proofs
 - Contradiction
 - Induction
- Common error: not checking/proving assumptions
 - "Jumps" in logic

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For Next Time

- Read first two pages of book's preface
- Read Chapter 1 of book