

Objectives

- Discussion of “The Status of P and NP”

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Review: Polynomial-Time Reduction

- **Purpose.** Classify problems according to *relative difficulty*.
- **Design algorithms.** If $X \leq_p Y$ and Y can be solved in polynomial-time, then X **can** also be solved in polynomial time.
- **Establish intractability.** If $X \leq_p Y$ and X cannot be solved in polynomial-time, then Y **cannot** be solved in polynomial time.
- **Establish equivalence.** If $X \leq_p Y$ and $Y \leq_p X$, we use notation $X \equiv_p Y$.

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Review: Basic Reduction Strategies

- Reduction by simple equivalence
- Reduction from special case to general case
- Reduction by encoding with gadgets

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Status of the P vs NP Problem

From Numbers

Charlie: Dad, uhm.. I've been working on a problem. P vs. NP. It can't be solved.

Alan: I think you knew that when you started.

Charlie: I could work on it forever. Constantly pushing forward, still never reaching an end.

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Status of the P vs NP Problem

- What is P? NP? What is the problem?

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NP

- Problems that no polytime algorithm has been found AND have not proven that no polytime algorithm exists
 - A little more ...
- Examples:

Name	Description
Hamiltonian circuit	Determine whether a given graph has a Hamiltonian circuit (a path that starts & ends at the same vertex and passes through all other vertices exactly once)
Traveling salesman	Find the shortest tour through n cities with known positive integer distances between them (each city once)
Graph coloring	Find a graph's chromatic number: smallest # of colors that need to be assigned to the graph's vertices so that no 2 adjacent vertices are assigned the same color.

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Common Feature

- Computationally difficult BUT checking if a proposed solution solves problem *can* be solved in polynomial time
- Example: easy to check if a proposed list of vertices is a Hamiltonian circuit for a given graph with n vertices

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Nondeterministic Algorithm

- Input: instance of a decision problem
- 1. Nondeterministic “guessing” stage: guess a solution to problem I
- 2. Deterministic “verification” stage: outputs yes if solution is a solution to the problem

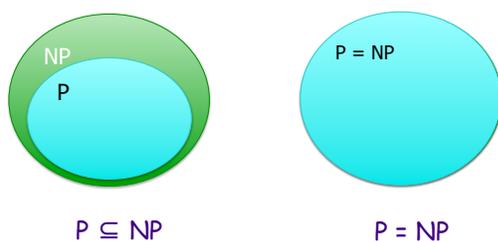
NP: A nondeterministic algorithm whose verification stage has a polynomial runtime.

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What We're Trying To Figure Out



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SAT-3

- *Satisfiability* is determining if the variables of a given Boolean formula can be assigned to make the formula evaluate to True
- Special case: 3 SAT
 - Formula is in conjunctive normal form
 - & between clauses (ors and nots)
 - Each clause contains at most 3 literals (variables)

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Status of the P vs NP Problem

- What are the consequences of NP Completeness?
- What if $P = NP$?
- How have people tried to prove $P \neq NP$?
 - Limitations? Still in progress?

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Final

- Usual rules
- Due next Friday, 5 p.m. (end of exams)
- Can use book, notes, my lecture notes, me (limited)
- No outside resources

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