

CSCI211: Problem Set 2

Due Friday, January 27

Points Possible: 25

1. **4 pts. (2.4)** Take the following list of functions and arrange them in ascending order of growth rate. That is, if function $g(n)$ immediately follows function $f(n)$ in your list, then it should be the case that $f(n)$ is $O(g(n))$.

$$\begin{array}{llll} f_1(n) = 2^{\sqrt{\log n}} & f_2(n) = 2^n & f_3(n) = n^{4/3} & f_4(n) = n(\log n)^3 \\ f_5(n) = n^{\log n} & f_6(n) = 2^{2^n} & f_7(n) = 2^{n^2} & \end{array}$$

2. **8 pts. Do chapter 2, exercise 6.** You may want to implement your algorithm to verify its correctness. Some Python code is available on the course web site to get you started.
3. **9 pts. (Levitan, 6.4.1)**
- (a) Construct a heap for the keys 1, 8, 6, 5, 3, 7, 4 by successive key insertions (top-down algorithm, i.e., Heapify-Up). You can draw your heap as an array or a tree. Tree will probably be easier. Show your work to ensure the possibility of partial credit.
 - (b) Construct a heap for the same keys (stored in an array) using the following bottom-up algorithm:

Algorithm 1 HeapBottomUp($H[1..n]$)

```
for  $i = \lfloor n/2 \rfloor$  downto 1 do
     $k = i$ 
     $v = H[k]$ 
     $heap = false$ 
    while not  $heap$  and  $2 * k \leq n$  do
         $j = 2 * k$ 
        if  $j < n$  then
            if  $H[j] > H[j + 1]$  then
                 $j = j + 1$ 
            end if
        end if
        if  $v \leq H[j]$  then
             $heap = true$ 
        else
             $H[k] = H[j]$ 
             $k = j$ 
        end if
    end while
     $H[k] = v$ 
end for
```

- (c) Compare the algorithms in terms of runtime and conceptual complexity.
 - (d) Will the successive-key insertions and bottom-up algorithms always yield the same heap for the same insertions?
4. 4 pts. (Levitan, 6.4.2) Design an algorithm for checking whether an array $H[1..n]$ is a heap and determine its runtime efficiency.