

## CSCI211: Problem Set 2

Due Friday, January 27

Points Possible: 25

1. **4 pts. (2.4)** Take the following list of functions and arrange them in ascending order of growth rate. That is, if function  $g(n)$  immediately follows function  $f(n)$  in your list, then it should be the case that  $f(n)$  is  $O(g(n))$ .

$$\begin{array}{llll} f_1(n) = 2^{\sqrt{\log n}} & f_2(n) = 2^n & f_3(n) = n^{4/3} & f_4(n) = n(\log n)^3 \\ f_5(n) = n^{\log n} & f_6(n) = 2^{2^n} & f_7(n) = 2^{n^2} & \end{array}$$

2. **8 pts. Do chapter 2, exercise 6.** You may want to implement your algorithm to verify its correctness. Some Python code is available on the course web site to get you started.
3. **9 pts. (Levitan, 6.4.1)**
- (a) Construct a heap for the keys 1, 8, 6, 5, 3, 7, 4 by successive key insertions (top-down algorithm, i.e., Heapify-Up). You can draw your heap as an array or a tree. Tree will probably be easier. Show your work to ensure the possibility of partial credit.
- (b) Construct a heap for the same keys (stored in an array) using the following bottom-up algorithm:

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**Algorithm 1** HeapBottomUp( $H[1..n]$ )

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for  $i = \lfloor n/2 \rfloor$  downto 1 do
   $k = i$ 
   $v = H[k]$ 
   $heap = false$ 
  while not  $heap$  and  $2 * k \leq n$  do
     $j = 2 * k$ 
    if  $j < n$  then
      if  $H[j] > H[j + 1]$  then
         $j = j + 1$ 
      end if
    end if
    if  $v \leq H[j]$  then
       $heap = true$ 
    else
       $H[k] = H[j]$ 
       $k = j$ 
    end if
  end while
   $H[k] = v$ 
end for
```

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- (c) Compare the algorithms in terms of runtime and conceptual complexity.
  - (d) Will the successive-key insertions and bottom-up algorithms always yield the same heap for the same insertions?
4. 4 pts. (Levitan, 6.4.2) Design an algorithm for checking whether an array  $H[1..n]$  is a heap and determine its runtime efficiency.