Objectives

- Review: Directed Graphs
- Topological Orderings of DAGs

Review

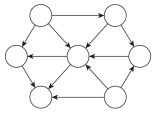
- What do we know about graphs? Directed graphs?

Review

- What do we know about graphs?
  - Space
  - Connectivity
    - BFS, DFS
  - Bipartite graphs
    - How to color?
    - When know not colorable?

Review: Representing Directed Graphs

- Edge (u, v) goes from node u to node v
- For each node, keep track of
  - Out edges (where links go)
  - In edges (from where links come in)

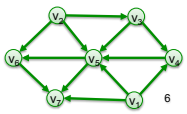


DAGS AND TOPOLOGICAL ORDERING

Directed Acyclic Graphs

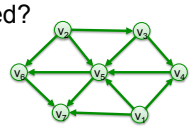
- Def. A ***DAG*** is a directed graph that contains no directed cycles.
- Example. Precedence constraints: edge (v_i, v_j) means v_i must precede v_j
  - Course prerequisite graph: course v_i must be taken before v_j
  - Compilation: module v_i must be compiled before v_j
  - Pipeline of computing jobs: output of job v_i needed to determine input of job v_j

a DAG:



Problem: Valid Ordering

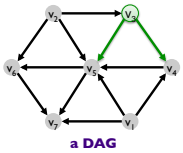
- Given a set of tasks with dependencies, what is a valid order in which the tasks could be performed?



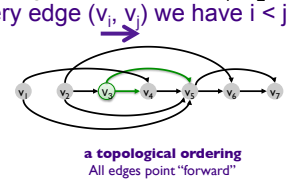
- Example: Getting dressed
  - What tasks are involved?
  - What tasks depend on other tasks?

Topological Ordering

- Problem: Given a set of tasks with dependencies, what is a valid order in which the tasks could be performed?
- Def. A ***topological order*** of a directed graph G = (V, E) is an ordering of its nodes as v_1, v_2, …, v_n so that for every edge (v_i, v_j) we have i < j.



a DAG



a topological ordering
All edges point "forward"

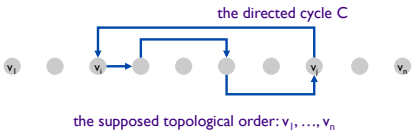
Coordinating labeling of nodes, but numbering is not known for just DAG

Towards a Solution

- Start by showing that if G has a topological order, then G is a DAG
- Eventually, we'll show the other direction: if G is a DAG, then G has a topological order

Directed Acyclic Graphs

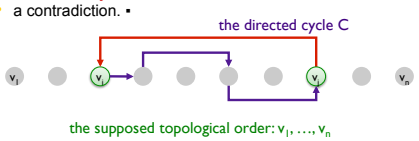
- Lemma. If G has a topological order, then G is a DAG.
- Proof plan: Try to show that G has a topological order even though G has a cycle



Why isn't this valid?

DAGs & Topological Orderings

- Lemma. If G has a topological order, then G is a DAG.
- Pf. (by contradiction)
  - Suppose that G has a topological order v_1, …, v_n and that G also has a directed cycle C.
  - Let v_i be the lowest-indexed node in C, and let v_j be the node on C just before v_i; thus (v_j, v_i) is an edge
  - By our choice of i (lowest-indexed node), i < j
  - Since (v_j, v_i) is an edge and v_1, …, v_n is a topological order, we must have j < i
    - a contradiction. ▪

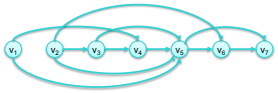


Directed Acyclic Graphs

- Does every DAG have a topological ordering?
  - If so, how do we compute one?

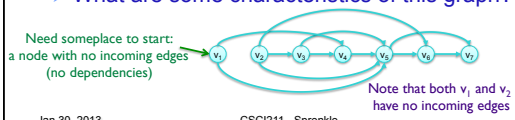
Directed Acyclic Graphs

- Does every DAG have a topological ordering?
  - If so, how do we compute one?
- What do we need to be able to create a topological ordering?
  - What are some characteristics of this graph?



Directed Acyclic Graphs

- Does every DAG have a topological ordering?
  - If so, how do we compute one?
- What do we need to be able to create a topological ordering?
  - What are some characteristics of this graph?



Towards a Topological Ordering

Goal: Find an algorithm for finding the TO
Idea: 1st node is one with no incoming edges

Do we know there is always a node with no incoming edges?

Towards a Topological Ordering

- Lemma. If G is a DAG, then G has a node with no incoming edges
  - This is our starting point of the topological ordering

How to prove?

Towards a Topological Ordering

- Lemma. If G is a DAG, then G has a node with no incoming edges
- Proof idea: Consider if there is no node without incoming edges
  - What contradiction are we looking for?

Towards a Topological Ordering

- Lemma. If G is a DAG, then G has a node with no incoming edges.
- Pf. (by contradiction)
  - Suppose that G is a DAG and every node has at least one incoming edge
  - Pick any node v, and follow edges backward from v.
    - Since v has at least one incoming edge (u, v), we can walk backward to u
  - Since u has at least one incoming edge (t, u), we can walk backward to t
  - Repeat until we visit a node, say w, twice
    - Has to happen at least by n+1 steps (Why?)
  - Let C denote the sequence of nodes encountered between successive visits to w. C is a cycle, which is a contradiction to G is a DAG ▪

Putting it all together: Creating a topological order

Ideas?

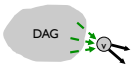
Topological Ordering Algorithm

```
Find a node v with no incoming edges
Order v first
Delete v from G
Recursively compute a topological ordering of G-{v}
      and append this order after v
```

How do we know this works?

Directed Acyclic Graphs

- Lemma. If G is a DAG, then G has a topological ordering.
- Pf. (by induction on n)
  - Base case: true if n = 1
  - Given DAG on n > 1 nodes, find a node v with no incoming edges
  - G - { v } is a DAG because deleting v cannot create cycles
  - By inductive hypothesis, G - { v } has a topological ordering
  - Place v first in topological ordering;
  - Append nodes of G - { v } in topological order.
    - valid since v has no incoming edges. ▪

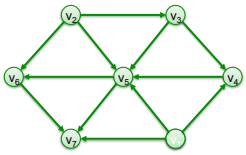


Topological Ordering Algorithm

- Lemma. If G is a DAG, then G has a topological ordering.
- Algorithm:

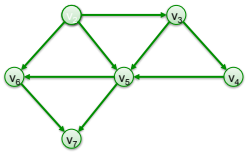
```
Find a node v with no incoming edges
Order v first
Delete v from G
Recursively compute a topological ordering of G-{v}
      and append this order after v
```


Topological Ordering Algorithm: Example



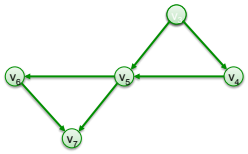
Topological order:

Topological Ordering Algorithm: Example



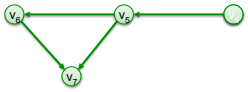
Topological order: v_1

Topological Ordering Algorithm: Example



Topological order: v_1, v_2

Topological Ordering Algorithm: Example



Topological order: v_1, v_2, v_3

Topological Ordering Algorithm: Example

Topological order: v_1, v_2, v_3, v_4

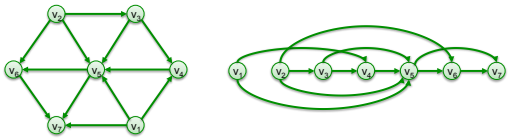
Topological Ordering Algorithm: Example

Topological order: v_1, v_2, v_3, v_4, v_5

Topological Ordering Algorithm: Example

Topological order: $v_1, v_2, v_3, v_4, v_5, v_6$

Topological Ordering Algorithm: Example



Topological order: $v_1, v_2, v_3, v_4, v_5, v_6, v_7$.

Topological Order Runtime

```
Find a node v with no incoming edges
Order v first
Delete v from G
Recursively compute a topological ordering of G-{v}
      and append this order after v
```

- Where are the costs?
- How would we implement?

Topological Order Runtime

```
Find a node v with no incoming edges
Order v first                                  O(n)
Delete v from G
Recursively compute a topological ordering of G-{v}   O(n)
      and append this order after v
```

- Find a node without incoming edges and delete it: **O(n)**
- Repeat on all nodes

Can we do better?

→ **$O(n^2)$**

Topological Sorting Algorithm: Running Time

- Theorem. Find a topological order in O(m + n) time
- Pf.
  - Maintain the following information:
    - count[w] = remaining number of incoming edges
    - S = set of remaining nodes with no incoming edges
  - Initialization: O(m + n) via single scan through graph
  - Algorithm:
    - Select a node v from S, remove v from S
    - Decrement count[w] for all edges from v to w
      - Add w to S if count[w] = 0

PS2 Feedback

- When providing algorithms, make sure your input is clear
  - Examples: what is the name of your heap? What does n represent?
  - `isHeap(H[1...n])`
- Analyze runtime of all algorithms created
- Consider implementing solutions
  - Catch errors when try different test cases
- I write notes on your algorithms so that I can understand what is happening

Looking Ahead

- Problem Set 3 due Friday
- Exam 1 handed out on Friday
  - Different rules from problem set
  - **No collaboration**
  - Can access your notes, book, my lectures
  - Can ask me questions, but I'm limited in how much help I can give
  - Next Wednesday: work session