

Objectives

- Analyzing proofs
- Introduction to problem solving
 - Our process, through an example

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Review

- What are our goals in solving problems?
- How do we show that our solutions are correct and efficient?
- What proof techniques did we discuss?

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Proof: All Horses Are The Same Color

- **Base case:** If there is only *one* horse, there is only one color.
- **Induction step:** Assume as induction hypothesis that within any set of n horses, there is only one color.
 - Look at any set of $n + 1$ horses
 - Label the horses: $1, 2, 3, \dots, n, n + 1$
 - Consider the sets $\{1, 2, 3, \dots, n\}$ and $\{2, 3, 4, \dots, n + 1\}$
 - Each is a set of only n horses, therefore within each there is only one color
 - Since the two sets overlap, there must be only one color among all $n + 1$ horses

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Where is the error in the proof?

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Error in Proof

- **Base case:** If there is only *one* horse, there is only one color.
- **Induction step:** Assume as induction hypothesis that within any set of n horses, there is only one color.
 - Look at any set of $n + 1$ horses
 - Number them: $1, 2, 3, \dots, n, n + 1$
 - Consider the sets $\{1, 2, 3, \dots, n\}$ and $\{2, 3, 4, \dots, n + 1\}$
 - Each is a set of only n horses, therefore within each there is only one color
 - Since the two sets overlap, there must be only one color among all $n + 1$ horses

Does not hold true when $n+1=2$

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Lesson: check assumptions within proof

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Proof Summary

- Need to **prove** conjectures
- Common types of proofs
 - Direct proofs
 - Contradiction
 - Induction
- Common error: not checking/proving assumptions
 - "Jumps" in logic

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Importance of Precision

- Frank and the toilet repair

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Process, through example

INTRODUCTION TO PROBLEM SOLVING

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Matching Residents to Hospitals

- **Goal:** Given a set of preferences among hospitals and medical school students, design a *self-reinforcing* admissions process.
- Applicant x and hospital y are *unstable* if
 - x prefers y to its assigned hospital
 - y prefers x to one of its admitted students
- **Stable assignment:** Assignment with no unstable pairs
 - No incentive for some pair of participants to undermine assignment by joint action
 - Unstable pair could each improve by swapping with current assignment
 - Self-reinforcing part

What details make this problem tricky?
What info do we need to solve problem?

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Stable Matching Problem

Simplified version of resident-matching problem

- **Goal:** Given n men and n women, find a "suitable" matching
 - Participants rank members of opposite sex
 - Each man lists women in order of preference from best to worst
 - Each woman lists men in order of preference from best to worst

	favorite ↓ 1 st	2 nd	least favorite ↓ 3 rd
Xavier	Amy	Bertha	Clare
Yancey	Bertha	Amy	Clare
Zeus	Amy	Bertha	Clare

Men's Preference Profile

	favorite ↓ 1 st	2 nd	least favorite ↓ 3 rd
Amy	Yancey	Xavier	Zeus
Bertha	Xavier	Yancey	Zeus
Clare	Xavier	Yancey	Zeus

Women's Preference Profile

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Stable Matching Goals

- **Perfect matching:** everyone is matched monogamously
 - Each man is paired with exactly one woman
 - Each woman is paired with exactly one man
- **Stability:** no incentive for some pair of participants to undermine assignment by joint action
 - An *unmatched* pair $m-w$ is *unstable* if man m and woman w prefer each other to current partners
 - Unstable pair $m-w$ could each improve by eloping
- **Stable matching:** perfect matching with no unstable pairs

Stable matching problem:

Given the preference lists of n men and n women, find a stable matching if one exists.

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Analyzing Stability

Instable: m prefers w to his woman; w prefers m to her man

- Is pairing X-C, Y-B, Z-A stable?

	favorite ↓ 1 st	2 nd	least favorite ↓ 3 rd
Xavier	Amy	Bertha	Clare
Yancey	Bertha	Amy	Clare
Zeus	Amy	Bertha	Clare

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Bertha	Xavier	Yancey	Zeus
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Women's Preference Profile

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Analyzing Stability

- Is pairing X-C, Y-B, Z-A stable?
- No. Bertha and Xavier prefer each other

	favorite ↓ 1 st	2 nd	least favorite ↓ 3 rd
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Stable Matching Problem

Instable: m prefers w to his woman; w prefers m to her man

- Is pairing X-A, Y-B, Z-C stable?
- Yes.

	favorite 1 st	2 nd	least favorite 3 rd		favorite 1 st	2 nd	least favorite 3 rd
Xavier	Amy	Bertha	Clare	Amy	Yancey	Xavier	Zeus
Yancey	Bertha	Amy	Clare	Bertha	Xavier	Yancey	Zeus
Zeus	Amy	Bertha	Clare	Clare	Xavier	Yancey	Zeus

Men's Preference Profile Women's Preference Profile

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Any Questions?

- What are you wondering about this problem/ its solution at this point?

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Any Questions?

- What are you wondering about this problem/ its solution at this point?
- Hopefully:
 - Is there a stable matching for every pair of preference lists?
 - If so, is there an algorithm to find the stable matching?
 - Can we be fair in the matching? (preferences)
 - Will the matching always be the same?

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Thoughts on Solving Problem

- What do we need to solve the problem?
- What do we know?
- Where should the state start?
- What are some initial ideas about approaches?

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Thoughts on Solving Problem

- Initially, no one is matched
- Pick an arbitrary man and have him match with his favorite woman.
 - Are we guaranteed that pair will be part of a stable matching?
- Should a woman accept her first offer? If not, what should she do?
- When are we done? Do we need to consider all combinations?

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Propose-And-Reject Algorithm

[Gale-Shapley 1962]

- Intuitive method that guarantees finding a stable matching

```

Initialize each person to be free
while (some man is free and hasn't proposed to every woman)
  Choose such a man m
  w = 1st woman on m's list to whom m has not yet proposed
  if w is free
    assign m and w to be engaged
  else if w prefers m to her fiancé m'
    assign m and w to be engaged and m' to be free
  else
    w rejects m
  
```

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Applying the Algorithm

	favorite			least favorite
	1 st	2 nd	3 rd	
Xavier	Amy	Bertha	Clare	
Yancey	Bertha	Amy	Clare	
Zeus	Amy	Bertha	Clare	

Men's Preference Profile

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Women's Preference Profile

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    w rejects m
  
```

Observations about the Algorithm

- What can we say about any woman's partner during the execution of the algorithm?
- How does a woman's state change over the execution of the algorithm?
- What can we say about a man's partner?

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Observations about the Algorithm

- What can we say about any woman's partner during the execution of the algorithm?
 - Observation 1. He gets "better" → she prefers him over her last partner
- How does a woman's state change over the execution of the algorithm?
 - Observation 2. Once a woman is matched, she never becomes unmatched; she only "trades up"
- What can we say about a man's partner?
 - Observation 3. She gets "worse"

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Proving Correctness

- Need to show
 - Algorithm terminates
 - Result is a perfect matching
 - Result is a stable matching



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Assignments

- Review Chapter 1
- Journal due next Tuesday
 - Preface
 - Chapter 1
 - Beginning of Chapter 2

I have a meeting at 4 p.m.,
so I won't be in my
office hours at that time.

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