

## Objectives

- Wrap-up Dijkstra's Algorithm
- Minimum Spanning Tree

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## Review: Greedy Algorithms and Dijkstra's Algorithm

- What are greedy algorithms?
- How are some strategies to prove that greedy algorithms are optimal?
- What was the greedy algorithm to find the shortest path in a weighted directed graph?

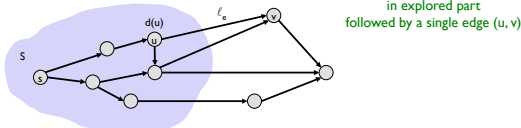
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## Dijkstra's Algorithm

1. Maintain a set of **explored nodes**  $S$ 
  - Keep the shortest path distance  $d(u)$  from  $s$  to  $u$
2. Initialize  $S = \{s\}$ ,  $d(s) = 0$ ,  $\forall u \neq s$ ,  $d(u) = \infty$
3. Repeatedly choose unexplored node  $v$  which minimizes  $\pi(v) = \min_{e = (u,v) : u \in S} d(u) + \ell_e$ ,
  - Add  $v$  to  $S$  and set  $d(v) = \pi(v)$ 
    - shortest path to some  $u$  in explored part followed by a single edge  $(u, v)$



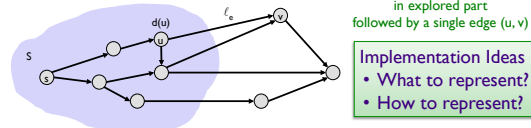
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## Dijkstra's Algorithm

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Implementation Ideas

- What to represent?
- How to represent?

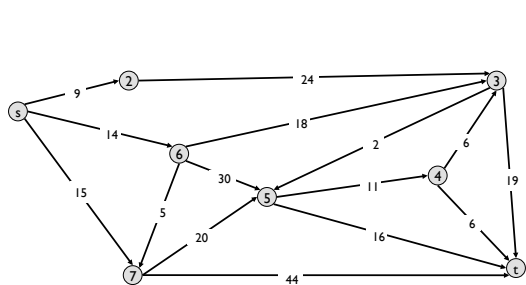
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## Dijkstra's Shortest Path Algorithm

- Find shortest path from  $s$  to  $t$  Give out handout



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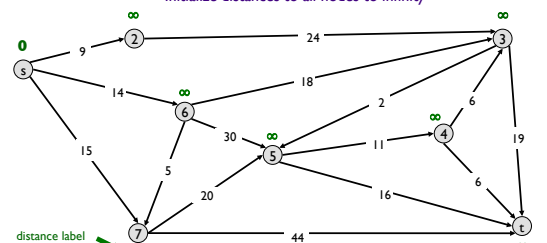
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## Dijkstra's Shortest Path Algorithm

$S = \{ \}$   
 $PQ = \{ s, 2, 3, 4, 5, 6, 7, t \}$

Initialize distances to all nodes to infinity



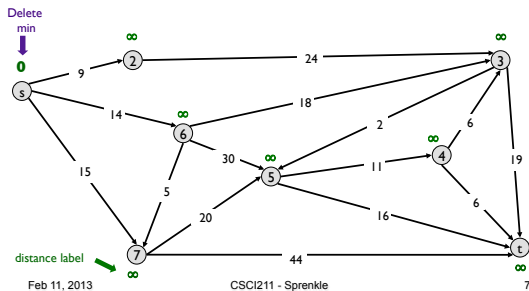
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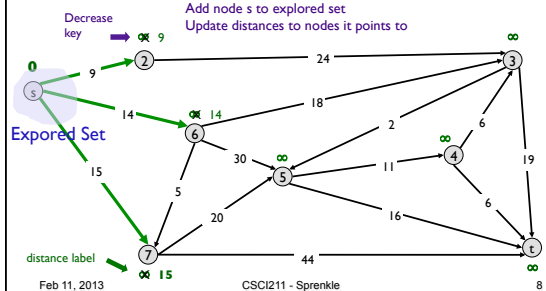
## Dijkstra's Shortest Path Algorithm

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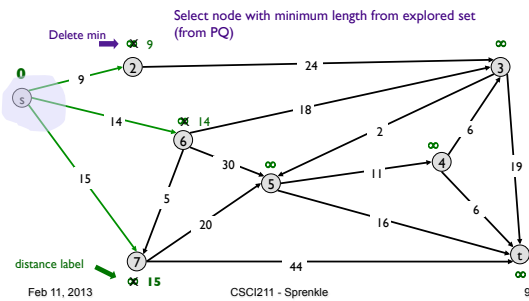
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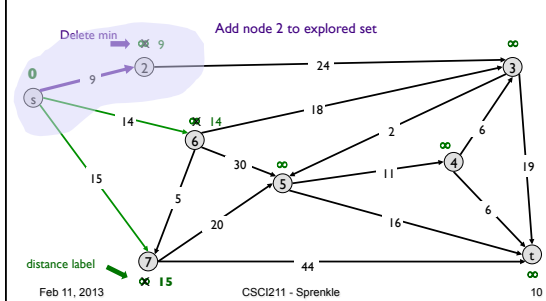
## Dijkstra's Shortest Path Algorithm

$S = \{s\}$   
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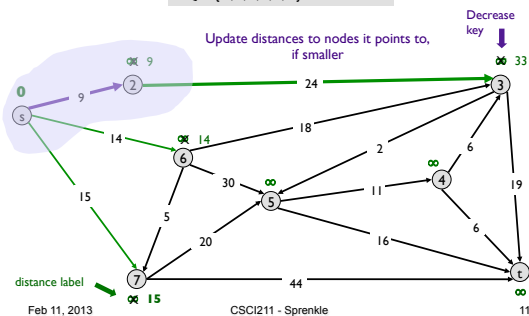
## Dijkstra's Shortest Path Algorithm

$S = \{s, 2\}$   
 $PQ = \{6, 7, 3, 4, 5, t\}$



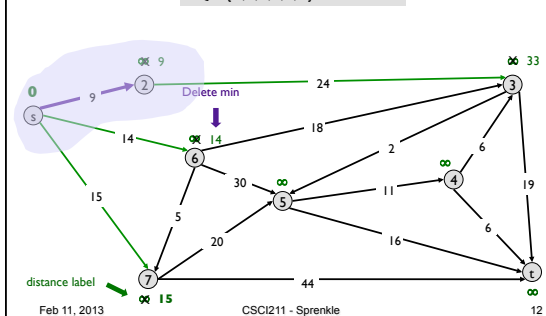
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$S = \{s, 2\}$   
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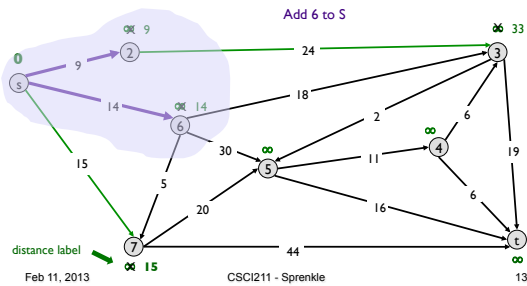


## Dijkstra's Shortest Path Algorithm

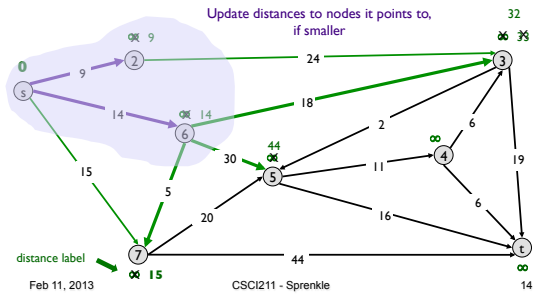
$S = \{s, 2\}$   
 $PQ = \{6, 7, 3, 4, 5, t\}$



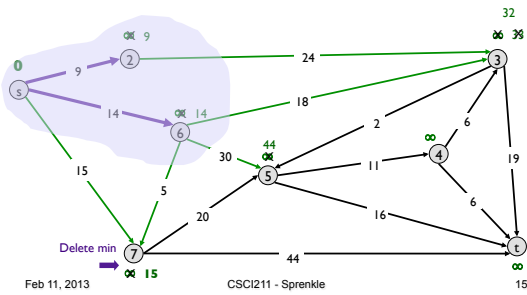
## Dijkstra's Shortest Path Algorithm

 $S = \{s, 2, 6\}$   
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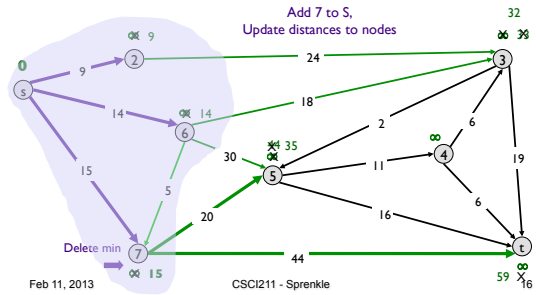
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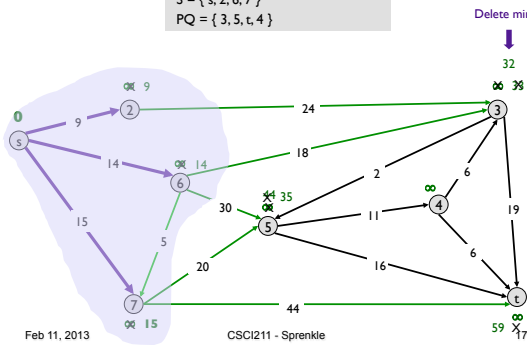
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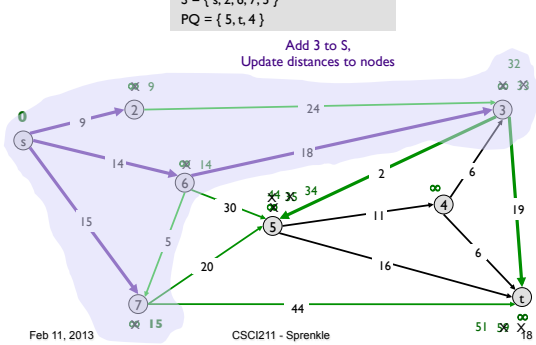
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 $S = \{s, 2, 6, 7\}$   
 $PQ = \{3, 5, t, 4\}$ 


## Dijkstra's Shortest Path Algorithm

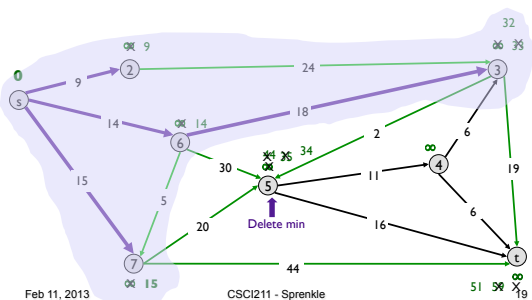
 $S = \{s, 2, 6, 7\}$   
 $PQ = \{3, 5, t, 4\}$ 


## Dijkstra's Shortest Path Algorithm

 $S = \{s, 2, 6, 7, 3\}$   
 $PQ = \{5, t, 4\}$ 


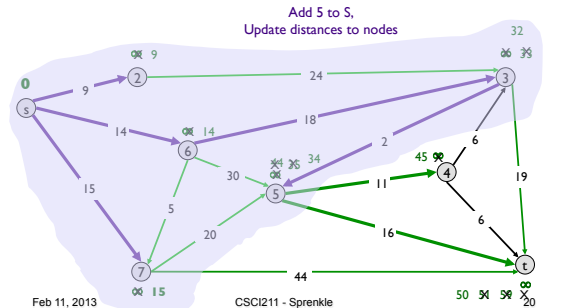
## Dijkstra's Shortest Path Algorithm

$S = \{s, 2, 6, 7, 3\}$   
 $PQ = \{5, t, 4\}$



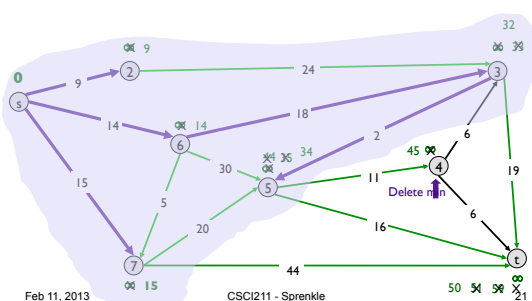
## Dijkstra's Shortest Path Algorithm

$S = \{s, 2, 6, 7, 3, 5\}$   
 $PQ = \{4, t\}$



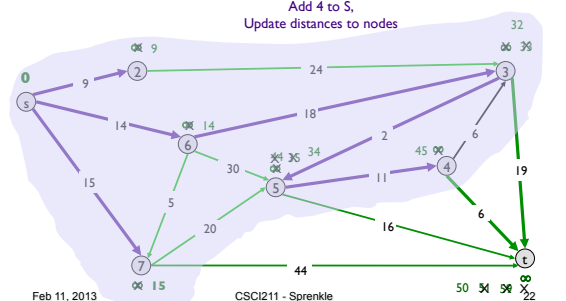
## Dijkstra's Shortest Path Algorithm

$S = \{s, 2, 6, 7, 3, 5\}$   
 $PQ = \{4, t\}$



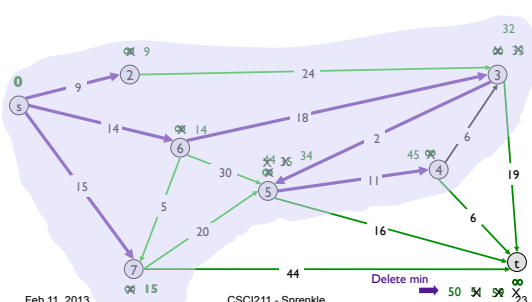
## Dijkstra's Shortest Path Algorithm

$S = \{s, 2, 6, 7, 3, 5, 4\}$   
 $PQ = \{t\}$



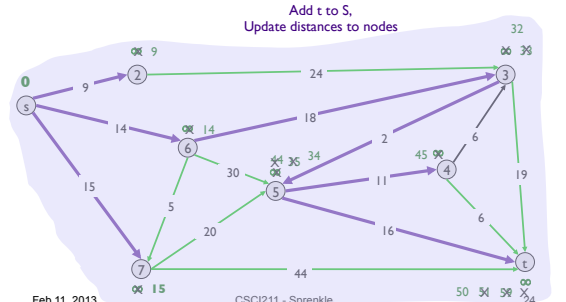
## Dijkstra's Shortest Path Algorithm

$S = \{s, 2, 6, 7, 3, 5, 4\}$   
 $PQ = \{t\}$



## Dijkstra's Shortest Path Algorithm

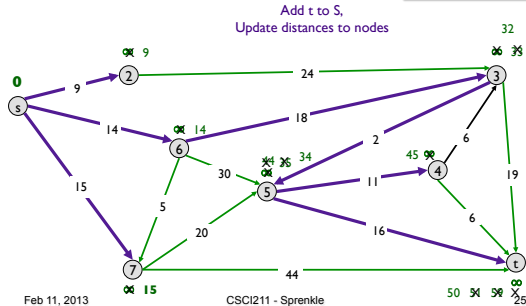
$S = \{s, 2, 6, 7, 3, 5, 4, t\}$   
 $PQ = \{\}$



## Dijkstra's Shortest Path Algorithm

$S = \{s, 2, 6, 7, 3, 5, 4, t\}$   
 $PQ = \{\}$

Why does Dijkstra's algorithm work?



## Dijkstra's Algorithm: Proof of Correctness

- **Invariant.** For each node  $u \in S$ ,  $d(u)$  is the length of the shortest  $s$ - $u$  path
- **Pf.** (by induction on  $|S|$ )
- **Base case:**  $|S|=1$  ...
- **Inductive hypothesis?**
- **Next step?**

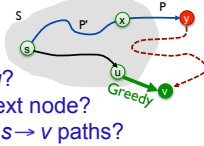
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## Dijkstra's Algorithm: Proof of Correctness

- **Prove:** For each node  $u \in S$ ,  $d(u)$  is the length of the shortest  $s$ - $u$  path
- **Pf.** (by induction on  $|S|$ )
- **Base case:** For  $|S| = 1$ ,  $S = \{s\}$ ;  $d(s) = 0$  ✓
- **Inductive hypothesis:** Assume true for  $|S| = k$ ,  $k \geq 1$ 
  - Grow  $|S|$  to  $k+1$
  - Greedy: Add node  $v$  by  $u \rightarrow v$
  - What do we know about  $s \rightarrow u$ ?
  - Why didn't we pick  $y$  as the next node?
  - What can we say about other  $s \rightarrow v$  paths?



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## Dijkstra's Algorithm: Proof of Correctness

- **Prove:** For each node  $u \in S$ ,  $d(u)$  is the length of the shortest  $s$ - $u$  path
- **Pf.** (by induction on  $|S|$ )
- **Inductive hypothesis:** Assume true for  $|S| = k$ ,  $k \geq 1$ 
  - Let  $v$  be the next node added to  $S$  by Greedy, and let  $u \rightarrow v$  be the chosen edge
  - The shortest  $s \rightarrow u$  path plus  $u \rightarrow v$  is an  $s \rightarrow v$  path of length  $\pi(v)$
  - Consider any  $s \rightarrow v$  path  $P$ . It's no shorter than  $\pi(v)$ .
  - Let  $x \rightarrow y$  be the first edge in  $P$  that leaves  $S$ , and let  $P'$  be the subpath to  $x$ .
  - $P$  is already too long as soon as it leaves  $S$ .

In terms of inequalities:

$$\ell(P) \geq \ell(P') + \ell(x, y) \geq d(x) + \ell(x, y) \geq \pi(y) \geq \pi(v)$$

nonnegative weights    inductive hypothesis    defn of  $\pi(y)$     Dijkstra chose  $v$  instead of  $y$

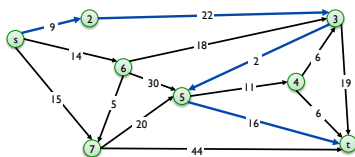
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## Discussion: Dijkstra's Algorithm

- Why does the algorithm break down if we allow negative weights/costs on edges?



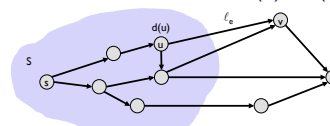
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## Dijkstra's Algorithm: Analysis

1. Maintain a set of explored nodes  $S$ 
  - Know the shortest path distance  $d(u)$  from  $s$  to  $u$
2. Initialize  $S = \{s\}$ ,  $d(s) = 0$ ,  $\forall u \neq s$ ,  $d(u) = \infty$
3. Repeatedly choose unexplored node  $v$  which minimizes  $\pi(v) = \min_{e = (u, v): u \in S} d(u) + \ell_e$ ,
  - Add  $v$  to  $S$  and set  $d(v) = \pi(v)$



shortest path to some  $u$  in explored part, followed by a single edge  $(u, v)$

Running time?  
Implementation?  
Data structures?

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## Dijkstra's Algorithm: Analysis

1. Maintain a set of explored nodes  $S$ 
  - Keep the shortest path distance  $d(u)$  from  $s$  to  $u$
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  - Add  $v$  to  $S$  and set  $d(v) = \pi(v)$

PQ Operation	RT of Op	# in Dijkstra
Insert		
ExtractMin		
ChangeKey		
IsEmpty		
Total		

- How long does each operation take?
- How many of each operation?

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## Dijkstra's Algorithm: Implementation

- For each unexplored node, explicitly maintain  $\pi(v) = \min_{e = (u,v) : u \in S} d(u) + \ell_e$ .
  - Next node to explore = node with minimum  $\pi(v)$ .
  - When exploring  $v$ , for each incident edge  $e = (v, w)$ , update  $\pi(w) = \min \{ \pi(w), \pi(v) + \ell_e \}$ .
- **Efficient implementation.** Maintain a priority queue of unexplored nodes, prioritized by  $\pi(v)$

PQ Operation	RT of Op	# in Dijkstra
Insert	$\log n$	$n$
ExtractMin	$\log n$	$n$
ChangeKey	$\log n$	$m$
IsEmpty	1	$n$
Total		$m \log n$

 **$O(m \log n)$** 

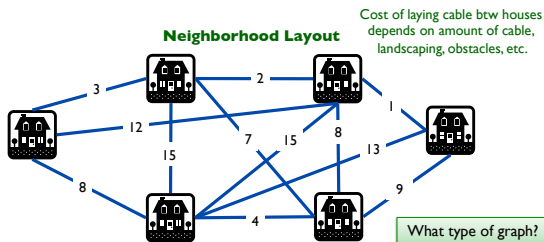
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## Laying Cable

- Comcast wants to lay cable in a neighborhood
  - Reach all houses
  - Least cost



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## Assignments

- Wiki due Tuesday
  - 3.4-3.6
  - Front matter of Chapter 4, 4.1
- Friday: PS4 Due

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