

## Objectives

- Analyzing proofs
- Introduction to problem solving
  - Our process, through an example

## Review

- What are our goals in solving problems?
- How do we show that our solutions are correct and efficient?
- What proof techniques did we discuss?

## Proof Summary

- Need to *prove* conjectures
- Common types of proofs
  - Direct proofs
  - Contradiction
  - Induction
- Common error: not checking/proving assumptions
  - “Jumps” in logic

## Proof: All Horses Are The Same Color

- **Base case:** If there is only *one* horse, there is only one color.
- **Induction step:** Assume as induction hypothesis that within any set of  $n$  horses, there is only one color.
  - Look at any set of  $n + 1$  horses
  - Label the horses:  $1, 2, 3, \dots, n, n + 1$
  - Consider the sets  $\{1, 2, 3, \dots, n\}$  and  $\{2, 3, 4, \dots, n + 1\}$
  - Each is a set of only  $n$  horses, therefore within each there is only one color
  - Since the two sets overlap, there must be only one color among all  $n + 1$  horses

Process, through example

## INTRODUCTION TO PROBLEM SOLVING

## Matching Residents to Hospitals

- **Goal:** Given a set of preferences among hospitals and medical school students, design a *self-reinforcing* admissions process.
- Applicant  $a$  and hospital  $h$  are *unstable* if
  - $a$  prefers  $h$  to its assigned hospital
  - $h$  prefers  $a$  to one of its admitted students
- **Stable assignment:** Assignment with no unstable pairs
  - No incentive for some pair of participants to undermine assignment by joint action
    - Unstable pair could each improve by swapping with current assignment
    - Self-

What details make this problem tricky?  
What info do we need to solve problem?

### Stable Matching Problem

*Simplified version of resident-matching problem*

- **Goal:** Given  $n$  men and  $n$  women, find a "suitable" matching
  - Participants rank members of opposite sex
  - Each man lists women in order of preference from best to worst
  - Each woman lists men in order of preference from best to worst

	favorite ↓ 1 <sup>st</sup>	2 <sup>nd</sup>	3 <sup>rd</sup>	least favorite ↓
Xavier	Amy	Bertha	Clare	
Yancey	Bertha	Amy	Clare	
Zeus	Amy	Bertha	Clare	

*Men's Preference Profile*

	favorite ↓ 1 <sup>st</sup>	2 <sup>nd</sup>	3 <sup>rd</sup>	least favorite ↓
Amy	Yancey	Xavier	Zeus	
Bertha	Xavier	Yancey	Zeus	
Clare	Xavier	Yancey	Zeus	

*Women's Preference Profile*

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### Stable Matching Goals

- **Perfect matching:** everyone is matched monogamously
  - Each man is paired with exactly one woman
  - Each woman is paired with exactly one man
- **Stability:** no incentive for some pair of participants to undermine assignment by joint action
  - An *unmatched* pair  $m-w$  is **unstable** if man  $m$  and woman  $w$  prefer each other to current partners
  - Unstable pair  $m-w$  could each improve by eloping
- **Stable matching:** perfect matching with no unstable pairs

#### Stable matching problem:

Given the preference lists of  $n$  men and  $n$  women, find a stable matching if one exists.

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### Analyzing Stability

**Instable:**  $m$  prefers  $w$  to his woman;  $w$  prefers  $m$  to her man

- Is pairing X-C, Y-B, Z-A stable?

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*Women's Preference Profile*

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### Analyzing Stability

- Is pairing X-C, Y-B, Z-A stable?
- No. Bertha and Xavier prefer each other

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### Stable Matching Problem

**Instable:**  $m$  prefers  $w$  to his woman;  $w$  prefers  $m$  to her man

- Is pairing X-A, Y-B, Z-C stable?
- Yes.

	favorite ↓ 1 <sup>st</sup>	2 <sup>nd</sup>	3 <sup>rd</sup>	least favorite ↓
Xavier	Amy	Bertha	Clare	
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*Women's Preference Profile*

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### Any Questions?

- What are you wondering about this problem/its solution at this point?

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### Any Questions?

- What are you wondering about this problem/its solution at this point?
- Hopefully:
  - Is there a stable matching for every pair of preference lists?
  - If so, is there an algorithm to find the stable matching?
  - Can we be fair in the matching? (preferences)
  - Will the matching always be the same?

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### Thoughts on Solving Problem

- What do we need to solve the problem?
- What do we know?
- Where should the state start?
- What are some initial ideas about approaches?

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### Thoughts on Solving Problem

- Initially, no one is matched
- Pick an arbitrary man and have him match with his favorite woman.
  - Are we guaranteed that pair will be part of a stable matching?
- Should a woman accept her first offer? If not, what should she do?
- When are we done? Do we need to consider all combinations?

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### Propose-And-Reject Algorithm

[Gale-Shapley 1962]

- Intuitive method that guarantees finding a stable matching

```

Initialize each person to be free
while (some man is free and hasn't proposed to every woman)
  Choose such a man m
  w = 1st woman on m's list to whom m has not yet proposed
  if w is free
    assign m and w to be engaged
  else if w prefers m to her fiancé m'
    assign m and w to be engaged and m' to be free
  else
    w rejects m
    
```

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### Applying the Algorithm

	favorite	1st	2nd	3rd	least favorite		favorite	1st	2nd	3rd	least favorite
Xavier		Amy	Bertha	Clare		Amy	Yancey	Xavier	Zeus		
Yancey		Bertha	Amy	Clare		Bertha	Xavier	Yancey	Zeus		
Zeus		Amy	Bertha	Clare		Clare	Xavier	Yancey	Zeus		

Men's Preference Profile                      Women's Preference Profile

```

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### Applying the Algorithm

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Men's Preference Profile                      Women's Preference Profile

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```

### Observations about the Algorithm

- What can we say about any woman’s partner during the execution of the algorithm?
- How does a woman’s state change over the execution of the algorithm?
- What can we say about a man’s partner?

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### Observations about the Algorithm

- What can we say about any woman’s partner during the execution of the algorithm?
  - Observation 1. He gets “better” → she prefers him over her last partner
- How does a woman’s state change over the execution of the algorithm?
  - Observation 2. Once a woman is matched, she never becomes unmatched; she only “trades up”
- What can we say about a man’s partner?
  - Observation 3. She gets “worse”

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### Proving Correctness

- Need to show
  - Algorithm terminates
  - Result is a perfect matching
  - Result is a stable matching



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### Propose-And-Reject Algorithm

[Gale-Shapley 1962]

Does algorithm terminate?

```

Initialize each person to be free
while (some man is free and hasn't proposed to every woman)
  Choose such a man m
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  if w is free
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    w rejects m
    
```

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### Proof of Correctness: Termination

- Claim. Algorithm terminates after at most  $n^2$  iterations of while loop.
  - Hint: How wouldn't the algorithm terminate?

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### Assignments

- Review Chapter 1
- Journal due Monday/Tuesday (because of MLK day)
  - Preface
  - Chapter 1
  - Beginning of Chapter 2

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